DEALING WITH ALIASING USING CONTRACTS

BEATING FORTRAN'S PERFORMANCE

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ALIASING

int f(int &a, float &b) {
 a = 2;
 b = 3;
 return a;
}

define i32 f(i32*, float*) {
 store i32 2, i32* %a
 store float 3, float* %b
 ret i32 2

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Some parameters might alias! Type based alias analysis

WHY DOES ALIASING MATTER?

LATENCY NUMBERS

L1 cache reference	0.5	ns						
Branch mispredict	5	ns						
L2 cache reference	7	ns	14x	L1	cache			
Mutex lock/unlock	25	ns						
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OPTIMIZATIONS

WHY DOES ALIASING MATTER?

LORE: LOop Repository for the Evaluation of compilers Numbers from P1296R0

	Loops	Sped up	Mean Speedup	Slowed	Mean slowdowr
GCC	1939	734 (38%)	2.39x	155 (8%)	0.766
ICC	1861	843 (45%)	2.59x	94 (5%)	0.61
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FORTRAN

- Procedure arguments and variables may not alias
- Inception when CPU time was expensive
- To convince people not to write in assembly...
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C++

No standard way (other than types) to give aliasing related hints.

NOT VECTORIZED

```
void f(int *a, int *b, const int& num) {
    for(int i = 0; i < num; ++i) {
        a[i] = b[i] * b[i] + 1;
    }
}</pre>
```

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Rings some bells?

Loop versioned, large unrolled code twice

Only the vectorized version

IS IT ALWAYS POSSIBLE TO UTILIZE THE TYPE BASED ALIASING RULES?

NOT VECTORIZED

void g(int *result, int **matrix, int height, int width) {
 for(int i = 0; i < height; ++i)
 for(int j = 0; j < width; ++j)
 result[i] += matrix[i][j];
</pre>

NOT VECTORIZED

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```

VECTORIZED

restrict

During each execution of a block in which a restricted pointer P is declared, if some object that is accessible through P (directly or indirectly) is modified, by any means, then all accesses to that object (both reads and writes) in that block must occur through P (directly or indirectly), otherwise the behavior is undefined.

LET'S JUST ADD RESTRICT TO C++?

How to annotate the code below?

```
void g(vector<int> &result, vector<vector<int>> &matrix) {
   for(int i = 0; i < matrix.size(); ++i)
      for(int j = 0; j < matrix[0].size(); ++j)
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```

What would

vector<int restrict>

or

vector<int> restrict

mean?

ADDING restrict TO C++

- Many failed attempts, lots of unanswered questions
- Should restrict change the overload sets?
- Should restrict participate in name mangling?
- restrict was never designed to work with the class abstraction
- How should restrict carried through templates?
- Members, lambda captures, unions, ...
- C2X, n2260, clarifying restrict

```
void f(int * restrict x, int * restrict y);
void g() {
    int x;
    f(&x, &x);
}
```

```
void f(int * restrict x, int * restrict y);
void g() {
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}
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Adding restrict to f makes it harder to use. It is now the caller's responsibility to ensure no aliasing is happening.

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Adding restrict to f makes it harder to use. It is now the caller's responsibility to ensure no aliasing is happening. Restrict is a precondition! Only if we had a way to describe preconditions in C++... Voted into C++20 in June (Rapperswil meeting)

CONTRACTS TO THE RESCUE? EXPLORING THE DESIGN SPACE

SIMPLE PRECONDITIONS

```
int f(int &a, int &b) [[expects axiom: &a != &b]] {
    a = 2;
    b = 3;
    return a;
}
```

- f(x, x); is undefined
- The precondition is documented
- We have two mitigations:
 - Runtime checks (with axiom removed)
 - Static analysis

SIMPLE PRECONDITIONS (LAMBDAS)

auto f = [](int &a, int &b) [[expects axiom: &a != &b]] {
 a = 2;
 b = 3;
 return a;
}

ARRAYS

int *merge(int *a, int *b, int num) [[expects: ???]];

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int *merge(int *a, int *b, int num)
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___disjoint(a, b, c, ..., num)?

ARRAYS

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Extend the language?

int *merge(int *a, int *b, int num)
 [[expects: __disjoint(a, b, num)]];

__disjoint(a, b, c, ..., num)?

int *merge(int *a, int *b, int num)
 [[expects: __distinct(a) && __distinct(b)]];

POSSIBLE IMPLEMENTATION FOR _____disjoint?

```
// From: P1296R0
template<typename T, typename U>
bool __disjoint(const T *pt, const U *pu, size_t n) {
    intptr_t bt = (intptr_t)pt,
        et = (intptr_t)(pt + n);
    intptr_t bu = (intptr_t)pu,
        eu = (intptr_t)(pu + n);
    return (et <= bu) || (eu <= bt);</pre>
```

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        eu = (intptr_t)(pu + n);
    return (et <= bu) || (eu <= bt);
}</pre>
```

Are we sure this is well defined? Compilers might want to have intrinsics instead.

USER DEFINED TYPES

int f(S a, S b)
[[expects:

[[expects: __disjoint(a.member, b.member)]];

USER DEFINED TYPES

int f(S a, S b)
 [[expects: disjoint(a.member, b.member)]];

int f(S a, S b)
 [[expects: ___disjoint(a.method(), b.method())]];

USER DEFINED TYPES

int f(S a, S b)
 [[expects: ___disjoint(a.member, b.member)]];

int f(S a, S b)
 [[expects: __disjoint(a.method(), b.method())]];

What if we need arguments? Use dummy symbols? Existentially or universally quantified?

int f(S a, S b)
 [[expects: __disjoint(a.method(???), b.method(???))]];

VIEWS TO THE RESCUE?

NON-ALIASING VIEW EXAMPLE

```
template <typename ... >
class unique_span {
    unique_span(...) [[expects: ???]];
    reference operator[](index_type idx) const
        [[ensures x: __distinct(x, this, idx)]];
};
```

f(unique_span(vec), unique_span(vec2));

BACK TO THE MATRIX EXAMPLE

```
void g(unique_span<int> result,
    vector<unique_span<int>> &matrix) {
    for(int i = 0; i < matrix.size(); ++i)
        for(int j = 0; j < matrix[0].size(); ++j)
            result[i] += matrix[i][j];
```

BACK TO THE MATRIX EXAMPLE

```
void g(unique_span<int> result,
    vector<unique_span<int>> &matrix) {
    for(int i = 0; i < matrix.size(); ++i)
        for(int j = 0; j < matrix[0].size(); ++j)
            result[i] += matrix[i][j];
```

Note that in real code we may want a multidimensional view or one dimensional matrix representation to avoid copying at the call site.

A NEW TYPE? ISN'T THAT HEAVY WEIGHT?

ARE THESE FUNCTIONS THE SAME?

double my_sqrt(double x) {
 return sqrt(x);

```
double my_sqrt(double x) {
    if (x < 0) return 0;
    return sqrt(x);
}</pre>
```

```
double my_sqrt(double x) {
    if (x < 0) throw ...;
    return sqrt(x);
}</pre>
```

ARE THESE FUNCTIONS THE SAME?

double my_sqrt(double x);

double my_sqrt(double x) [[expects: x >= 0]];

double my_sqrt(double x) [[expects: x >= 0]]
 [[ensures ret: ret >= 0]];

ARE THESE TYPES THE SAME?

unique_span<int>

span<int>

Exercise: how different are these types?

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- These can be vocabulary types
- We should use such classes sparingly, as they impose burden on the caller
- Those methods/functions are now screaming that they are special and error prone
- We can do overloads!

WE ALREADY HAVE TO REASON ABOUT ALIASING

- std::copy*
- memcpy vs memmove
- We would get mitigations for existing UB!

RELATED WORK

- **p0856r0**: Restrict as a library feature
- n3635, n4150: Annotating alias sets
- **P1296R0**: Very similar design, cooperating with the authors
- The malloc attribute of GCC, noalias attribute of Clang
- All major compilers has restrict like features as extensions
- IBM XL's #pragma disjoint

P1296R0

- std::disjoint only
- Discussed at San Diego meeting
- In ealry stages, no way to get into C++20

ALIAS SETS

```
void * [[alias_set()]] malloc(size_t);
int * [[alias_set(Foo)]] p1 = ...;
int * [[alias_set(Bar), alias_set(Baz)]] p2 = ...;
int * p3 = ...;
```

THANKS FOR YOUR ATTENTION!